

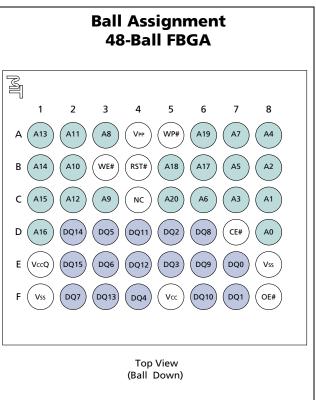
FLASH MEMORY

MT28F322P3

LOW VOLTAGE, EXTENDED TEMPERATURE

Features

- Flexible dual-bank architecture Support for true concurrent operation with zero latency Read bank *a* during program bank *b* and vice versa Read bank *a* during erase bank *b* and vice versa
- Basic configuration: Seventy-one erasable blocks Bank *a* (8Mb for data storage) Bank *b* (24Mb for program storage)
- VCC, VCCQ, VPP voltages
 2.7V (MIN), 3.3V (MAX) VCC
 2.2V (MIN), 3.3V (MAX) VCCQ
 3.0V (TYP) VPP (in-system PROGRAM/ERASE)
 12V ±5% (HV) VPP tolerant (factory programming compatibility)
- Random access time: 70ns @ 2.7V VCC
- Page Mode read access Eight-word page Interpage read access: 70ns @ 2.7V Intrapage read access: 30ns @ 2.7V
- Low power consumption (VCC = 3.3V) Asynchronous/interpage READ < 15mA Intrapage READ < 7mA WRITE < 20mA (MAX) ERASE < 25mA (MAX) Standby < 15μA (TYP), 50μA (MAX) @ 3.3V Automatic power save (APS) feature
- Enhanced write and erase suspend options ERASE-SUSPEND-to-READ within same bank PROGRAM-SUSPEND-to-READ within same bank ERASE-SUSPEND-to-PROGRAM within same bank
- Dual 64-bit chip protection registers for security purposes
- Cross-compatible command support Extended command set Common flash interface
- PROGRAM/ERASE cycle 100,000 WRITE/ERASE cycles per block
- Fast programming algorithm
- $VPP = 12V \pm 5\%$



NOTE: See page 7 for Ball Description Table. See page 39 for mechanical drawing.

OPTIONS

MARKING

•	Timing	
	70ns access	-70
	80ns access	-80
٠	Boot Block Configuration	
	Тор	Т
	Bottom	В
٠	Package	
	48-ball FBGA (6 x 8 ball grid)	FJ
•	Operating Temperature Range	
	Extended (-40°C to +85°C)	ET
	Part Number Example:	

MT28F322P3FJ-70 BET

09005aef808cfe29 MT28F322P3FJ_F.fm - Rev. F 7/03 EN



General Description

The MT28F322P3 is a high-performance, high-density, nonvolatile memory solution that can significantly improve system performance. This new architecture features a two-memory-bank configuration that supports background operation with no latency.

A high-performance bus interface allows a fast page mode data transfer; a conventional asynchronous bus interface is provided as well.

The MT28F322P3 allows soft protection for blocks, as read only, by configuring soft protection registers with dedicated command sequences. For security purposes, two 64-bit chip protection registers are provided.

The embedded WORD WRITE and BLOCK ERASE functions are fully automated by an on-chip write state machine (WSM). Two on-chip status registers, one for each of the two memory partitions, can be used to monitor the WSM status and to determine the progress of the program/erase \k.

The erase/program suspend functionality allows compatibility with existing EEPROM emulation software packages.

The device is manufactured using 0.18µm process technology.

2 MEG x 16 ASYNC/PAGE FLASH MEMORY

Refer to Micron's Web site www.micron.com/flash for the latest data sheet.

Architecture and Memory Organization

The MT28F322P3 Flash device contains two separate banks of memory (bank a and bank b) for simultaneous READ and WRITE operations.

The MT28F322P3 Flash memory is available in the following bank segmentation configuration:

- Bank *a* comprises one-fourth of the memory and contains 8 x 4K-word parameter blocks and 15 x 32K-word blocks.
- Bank *b* represents three-fourths of the memory, is equally sectored, and contains 48 x 32K-word blocks.

Figure 2 and Figure 3 show the bottom and top memory organizations.

Device Marking

Due to the size of the package, Micron's standard part number is not printed on the top of each device. Instead, an abbreviated device mark comprised of a five-digit alphanumeric code is used. The abbreviated device marks are cross referenced to Micron part numbers in Table 1.

PART NUMBER	PRODUCT MARKING	SAMPLE MARKING	MECHANICAL SAMPLE MARKING
MT28F322P3FJ-70 BET	FW816	FX816	FY816
MT28F322P3FJ-70 TET	FW817	FX817	FY817
MT28F322P3FJ-80 BET	FW814	FX814	FY814
MT28F322P3FJ-80 TET	FW815	FX815	FY815

 Table 1:
 Cross Reference for Abbreviated Device Marks





Part Numbering Information

Micron's low-power devices are available with several different combinations of features (see Figure 1). Valid combinations of features and their corresponding part numbers are listed in Table 2.

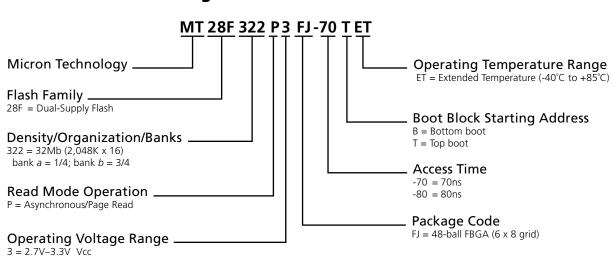


Table 2: Valid Part Number Combinations

PART NUMBER	ACCESS TIME (ns)	BOOT BLOCK STARTING ADDRESS	OPERATING TEMPERATURE RANGE
MT28F322P3FJ-70 BET	70	Bottom	-40°C to +85°C
MT28F322P3FJ-70 TET	70	Тор	-40°C to +85°C
MT28F322P3FJ-80 BET	80	Bottom	-40°C to +85°C
MT28F322P3FJ-80 TET	80	Тор	-40°C to +85°C

Figure 1: Part Number Chart



Functional Block Diagram

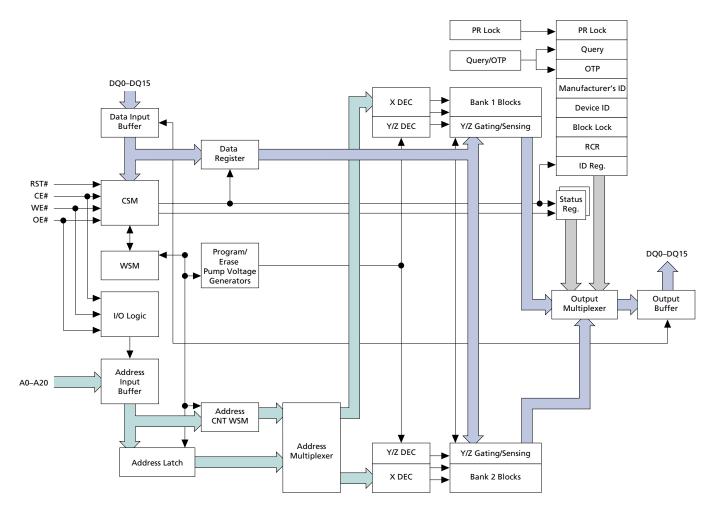




Figure 2: Bottom Boot Block Device

BANK <i>b</i> = 24Mb					
BLOCK	BLOCK SIZE (K-bytes/K-words)	ADDRESS RANGE (x16)			
70	64/32	1F8000h-1FFFFFh			
69	64/32	1F0000h-1F7FFFh			
68	64/32	1E8000h-1EFFFFh			
67	64/32	1E0000h-1E7FFFh			
66	64/32	1D8000h-1DFFFFh			
65	64/32	1D0000h-1D7FFFh			
64	64/32	1C8000h-1CFFFFh			
63	64/32	1C0000h-1C7FFFh			
62	64/32	1B8000h-1BFFFFh			
61	64/32	1B0000h-1B7FFFh			
60	64/32	1A8000h-1AFFFFh			
59	64/32	1A0000h-1A7FFFh			
58	64/32	198000h-19FFFFh			
57	64/32	190000h-197FFFh			
56	64/32	188000h-18FFFFh			
55	64/32	180000h-187FFFh			
54	64/32	178000h-17FFFFh			
53	64/32	170000h-177FFFh			
52	64/32	168000h-16FFFFh			
51	64/32	160000h-167FFFh			
50	64/32	158000h-15FFFFh			
49	64/32	150000h-157FFFh			
48	64/32	148000h-14FFFFh			
47	64/32	140000h-147FFFh			
46	64/32	138000h-13FFFFh			
45	64/32	130000h-137FFFh			
44	64/32	128000h-12FFFFh			
43	64/32	120000h-127FFFh			
42	64/32	118000h-11FFFFh			
41	64/32	110000h-117FFFh			
40	64/32	108000h-10FFFFh			
39	64/32	100000h-107FFFh			
38	64/32	0F8000h-0FFFFFh			
37	64/32	0F0000h-0F7FFFh			
36	64/32	0E8000h-0EFFFFh			
35	64/32	0E0000h-0E7FFFh			
34	64/32	0D8000h-0DFFFFh			
33	64/32	0D0000h-0D7FFFh			
32	64/32	0C8000h-0CFFFFh			
31	64/32	0C0000h-0C7FFFh			
30	64/32	0B8000h-0BFFFFh			
29	64/32	0B0000h-0B7FFFh			
23	64/32	0A8000h-0AFFFFh			
28	64/32	0A0000h-0A7FFFh			
26	64/32	098000h-097FFFh			
20	64/32	090000h-097FFFh			
23	64/32	088000h-087FFFh			
24	64/32	088000h-087FFFh			
23	04/32	00000011-08/FFF1)			

	BANK a = 8Mb					
BLOCK	BLOCK SIZE (K-bytes/K-words)	ADDRESS RANGE (x16)				
22	64/32	078000h-07FFFFh				
21	64/32	070000h-077FFFh				
20	64/32	068000h-067FFFh				
19	64/32	060000h-067FFFh				
18	64/32	058000h-05FFFFh				
17	64/32	050000h-057FFFh				
16	64/32	048000h-04FFFFh				
15	64/32	040000h-047FFFh				
14	64/32	038000h-03FFFFh				
13	64/32	030000h-037FFFh				
12	64/32	028000h-02FFFFh				
11	64/32	020000h-027FFFh				
10	64/32	018000h-01FFFFh				
9	64/32	010000h-017FFFh				
8	64/32	008000h-00FFFFh				
7	8/4	007000h-007FFFh				
6	8/4	006000h-006FFFh				
5	8/4	005000h-005FFFh				
4	8/4	004000h-004FFFh				
3	8/4	003000h-003FFFh				
2	8/4	002000h-002FFFh				
1	8/4	001000h-001FFFh				
0	8/4	000000h-000FFFh				



Figure 3: Top Boot Block Device

	BANK a = 8Mb					
BLOCK	BLOCK SIZE (K-bytes/K-words)	ADDRESS RANGE (x16)				
70	8/4	1FF000h-1FFFFFh				
69	8/4	1FE000h-1FEFFFh				
68	8/4	1FD000h-1FDFFFh				
67	8/4	1FC000h-1FCFFFh				
66	8/4	1FB000h-1FBFFFh				
65	8/4	1FA000h-1FAFFFh				
64	8/4	1F9000h-1F9FFFh				
63	8/4	1F8000h-1F8FFFh				
62	64/32	1F0000h-1F7FFFh				
61	64/32	1E8000h-1EFFFFh				
60	64/32	1E0000h-1E7FFFh				
59	64/32	1D8000h-1DFFFFh				
58	64/32	1D0000h-1D7FFFh				
57	64/32	1C8000h-1CFFFFh				
56	64/32	1C0000h-1C7FFFh				
55	64/32	1B8000h-1BFFFFh				
54	64/32	1B0000h-1B7FFFh				
53	64/32	1A8000h-1AFFFFh				
52	64/32	1A0000h-1A7FFFh				
51	64/32	198000h-19FFFFh				
50	64/32	190000h-197FFFh				
49	64/32	188000h-18FFFFh				
48	64/32	180000h-187FFFh				

BANK <i>b</i> = 24Mb					
BLOCK	BLOCK SIZE (K-bytes/K-words)	ADDRESS RANGE (x16)			
47	64/32	178000h-17FFFFh			
46	64/32	170000h-177FFFh			
45	64/32	168000h-16FFFFh			
44	64/32	160000h-167FFFh			
43	64/32	158000h-15FFFFh			
42	64/32	150000h-157FFFh			
41	64/32	148000h-14FFFFh			
40	64/32	140000h-147FFFh			
39	64/32	138000h-13FFFFh			
38	64/32	130000h-137FFFh			
37	64/32	128000h-12FFFFh			
36	64/32	120000h-127FFFh			
35	64/32	118000h-11FFFFh			
34	64/32	110000h-117FFFh			
33	64/32	108000h-10FFFFh			
32	64/32	100000h-107FFFh			
31	64/32	0F8000h-0FFFFFh			
30	64/32	0F0000h-0F7FFFh			
29	64/32	0E8000h-0EFFFFh			
28	64/32	0E0000h-0E7FFFh			
27	64/32	0D8000h-0DFFFFh			
26	64/32	0D0000h-0D7FFFh			
25	64/32	0C8000h-0CFFFFh			
24	64/32	0C0000h-0C7FFFh			
23	64/32	0B8000h-0BFFFFh			
22	64/32	0B0000h-0B7FFFh			
21	64/32	0A8000h-0AFFFFh			
20	64/32	0A0000h-0A7FFFh			
19	64/32	098000h-09FFFFh			
18	64/32	090000h-097FFFh			
17	64/32	088000h-08FFFFh			
16	64/32	080000h-087FFFh			
15	64/32	078000h-07FFFFh			
14	64/32	070000h-077FFFh			
13	64/32	068000h-06FFFFh			
12	64/32	060000h-067FFFh			
11	64/32	058000h-05FFFFh			
10	64/32	050000h-057FFFh			
9	64/32	048000h-04FFFFh			
8	64/32	040000h-047FFFh			
7	64/32	038000h-03FFFFh			
6	64/32	030000h-037FFFh			
5	64/32	028000h-02FFFFh			
4	64/32	020000h-027FFFh			
3	64/32	018000h-01FFFFh			
2	64/32	010000h-017FFFh			
1	64/32	008000h-00FFFFh			
0	64/32	000000h-007FFFh			
	1				



Ball Descriptions

48-BALL FBGA NUMBERS	SYMBOL	ТҮРЕ	DESCRIPTION
D8, C8, B8, C7, A8, B7, C6, A7, A3, C3, B2, A2, C2, A1, B1, C1, D1, B6, B5, A6, C5	A0–A20	Input	Address Inputs: Inputs for the address during READ and WRITE operations. Addresses are internally latched during READ and WRITE cycles.
D7	CE#	Input	Chip Enable: Activates the device when LOW. When CE# is HIGH, the device is disabled and goes into standby power mode.
F8	OE#	Input	Output Enable: Enables the output buffer when LOW. When OE# is HIGH, the output buffers are disabled.
В3	WE#	Input	Write Enable: Determines if a given cycle is a WRITE cycle. if WE# is LOW, the cycle is either a WRITE to the command state machine (CSM) or to the memory array.
B4	RST#	Input	Reset: When RST# is a logic LOW#, the device is in reset mode, which drives the outputs to High-Z and resets the write state machine (WSM). When RST# is at logic HIGH, the device is in standard operation. When RST# transitions from logic LOW to logic HIGH, the device resets all blocks to locked and defaults to the read array mode.
A5	WP#	Input	Write Protect: Controls the lock down function of the flexible locking feature.
A4	Vpp	Input	Program/Erase Enable: [1.8V–3.3V] Operates as input at logic levels to control complete device protection. Provides factory programming compatibility when driven to 11.4V–12.6V.
E7, F7, D5, E5, F4, D3, E3, F2, D6, E6, F6, D4, E4, F3, D2, E2	DQ0-DQ15	Input/Output	Data Inputs/Outputs: Input array data on the second CE# and WE# cycle during PROGRAM command. Input commands to the command user interface when CE# and WE# are active. DQ0–DQ15 output data when CE# and OE# are active.
E8, F1	Vss	Supply	Do not float any ground ball.
F5	Vcc	Supply	Device Power Supply: [2.7V–3.3V] Supplies power for device operation.
E1	VccQ	Supply	I/O Power Supply: [2.2V–3.3V] Supplies power for input/output buffers.
C4	NC	-	Internally not connected.



Command State Machine (CSM)

Commands are issued to the command state machine (CSM) using standard microprocessor write timings. The CSM acts as an interface between external microprocessors and the internal WSM. The available commands are listed in Table 3, their definitions are given in Table 4, and their descriptions in Table 5. Program and erase algorithms are automated by an on-chip WSM. For more specific information about the CSM transition states, see Micron technical note TN-28-33, Command State Machine Description and Command Definition.

Once a valid PROGRAM/ERASE command is entered, the WSM executes the appropriate algorithm, which generates the necessary timing signals to control the device internally and accomplish the requested operation. A command is valid only if the exact sequence of WRITEs is completed. After the WSM completes its task, the WSM status bit (SR7) (see Table 8) is set to a logic HIGH level (1), allowing the CSM to respond to the full command set again.

Operations

Device operations are selected by entering a standard JEDEC 8-bit command code with conventional microprocessor timings into an on-chip CSM through I/Os DQ0–DQ7. The number of bus cycles required to activate a command is typically one or two. The first operation is always a WRITE. Control signals CE# and WE# must be at a logic LOW level (VIL), and OE# and RST# must be at logic HIGH (VIH). The second operation, when needed, can be a WRITE or a READ depending upon the command. During a READ operation, control signals CE# and OE# must be at a logic LOW level (VIL), and WE# and RST# must be at logic HIGH (VIH).

Table 6 shows the bus operations for all the modes: write, read, reset, standby, and output disable.

When the device is powered up, internal reset circuitry initializes the chip to a read array mode of operation. Changing the mode of operation requires that a command code be entered into the CSM. For each one of the two memory partitions, an on-chip status register is available. These two registers allow the progress of the various operations that can take place on a memory bank to be monitored. One of the two status registers is interrogated by entering a READ STATUS REGISTER command onto the CSM (cycle 1), specifying an address within the memory partition boundary, and reading the register data on I/Os DQ0–DQ7 (cycle 2). Status register bits SR0–SR7 correspond to DQ0– DQ7 (see Table 8).

Command Definition

Once a specific command code has been entered, the WSM executes an internal algorithm, generating the necessary timing signals to program, erase, and verify data. See Table 4 for the CSM command definitions and data for each of the bus cycles.



Status Register

The status register allows the user to determine whether the state of a PROGRAM/ERASE operation is pending or complete. The status register is monitored by toggling OE# and CE# and reading the resulting status code on I/Os DQ0–DQ7. The high-order I/Os

(DQ8–DQ15) are set to 00h internally, so only the low order I/Os (DQ0–DQ7) need to be interpreted. Address lines select the status register pertinent to the selected memory partition.

Register data is updated and latched on the falling edge of OE# or CE#, whichever occurs last. Latching the data prevents errors from occurring if the register input changes during a status register read.

The status register provides the internal state of the WSM to the external microprocessor. During periods when the WSM is active, the status register can be polled to determine the WSM status. Table 8 defines the status register bits.

After monitoring the status register during a PRO-GRAM/ERASE operation, the data appearing on DQ0– DQ7 remains as status register data until a new command is issued to the CSM. To return the device to other modes of operation, a new command must be issued to the CSM.

Command State Machine Operations

The CSM decodes instructions for read array, read protection configuration register, read query, read status register, clear status register, program, erase, erase suspend, erase resume, program suspend, program resume, lock block, unlock block, and lock down block, chip protection program, and set read configuration register. The 8-bit command code is input to the device on DQ0–DQ7 (see Table 3 for CSM codes, Table 4 for command definitions, and Table 6 for CSM transition states). During a PROGRAM or ERASE cycle, the CSM informs the WSM that a PROGRAM or ERASE cycle has been requested.

During a PROGRAM cycle, the WSM controls the program sequences and the CSM responds to a PRO-GRAM SUSPEND command only.

During an ERASE cycle, the CSM responds to an ERASE SUSPEND command only. When the WSM has completed its task, the WSM status bit (SR7) is set to a logic HIGH level and the CSM responds to the full command set. The CSM stays in the current command state until the microprocessor issues another command.

The WSM successfully initiates an ERASE or PRO-GRAM operation only when VPP is within its correct voltage range.

COMMAND DQ0-DQ7	CODE ON DEVICE MODE
40h/10h	Program setup/alternate program setup
20h	Block erase setup
30h	Fast programming algorithm setup
50h	Clear status register
60h	Protection configuration setup
70h	Read status register
90h	Read protection configuration register
98h	Read query
B0h	Program/erase suspend
C0h	Protection register program/lock
D0h	Program/erase resume – erase confirm
FFh	Read array
01h	Lock Block
2Fh	Lock Down

Table 3: Command State Machine Codes for Device Mode Selection



Table 4: Command Definitions

	FIRST BUS CYCLE			SECOND BUS CYCLE		
COMMAND	OPERATION	ADDRESS	DATA	OPERATION	ADDRESS	DATA
READ ARRAY	WRITE	WA	FFh			
READ PROTECTION CONFIGURATION REGISTER	WRITE	IA	90h	READ	IA	ID
READ STATUS REGISTER	WRITE	BA	70h	READ	Х	SRD
CLEAR STATUS REGISTER	WRITE	BA	50h			
READ QUERY	WRITE	QA	98h	READ	QA	QD
BLOCK ERASE SETUP	WRITE	BA	20h	WRITE	BA	D0h
PROGRAM SETUP/ALTERNATE PROGRAM SETUP	WRITE	WA	40h/10h	WRITE	WA	WD
FAST PROGRAMMING ALGORITHM SETUP	WRITE	WA	30h	WRITE	WA	D0h
PROGRAM/ERASE SUSPEND	WRITE	BA	B0h			
PROGRAM/ERASE RESUME - ERASE CONFIRM	WRITE	BA	D0h			
LOCK BLOCK	WRITE	BA	60h	WRITE	BA	01h
UNLOCK BLOCK	WRITE	BA	60h	WRITE	BA	D0h
LOCK DOWN BLOCK	WRITE	BA	60h	WRITE	BA	2Fh
PROTECTION REGISTER PROGRAM	WRITE	PA	C0h	WRITE	PA	PD
PROTECTION REGISTER LOCK	WRITE	LPA	C0h	WRITE	LPA	FFFDh

NOTE: 1.BA:Address within the block IA: Identification code address ID: Identification code data LPA: Lock protection register address PA: Protection register address PD: Data to be written at location PA QA: Query code address QD: Query code data

SRD: Data read from the status register WA: Word address of memory location to be written, or read WD: Data to be written at the location WA

X: "Don't Care"



Table 5: Command Descriptions

CODE	DEVICE MODE	BUS CYCLE	DESCRIPTION
10h	Alt. Program Setup	First	Operates the same as a PROGRAM SETUP command.
20h	Erase Setup	First	Prepares the CSM for an ERASE CONFIRM command. If the next command is not an ERASE CONFIRM command, the command will be ignored, and the bank will go to the read array mode and wait for another command.
30h	FPA Setup	First	Prepares the CSM for an FPA CONFIRM command.
40h	Program Setup	First	A two-cycle command: The first cycle prepares for a PROGRAM operation, and the second cycle latches addresses and data and initiates the WSM to execute the program algorithm. The Flash device outputs status register data on the falling edge of OE# or CE#, whichever occurs first.
50h	Clear Status Register	First	The WSM can set the block lock status (SR3), program status (SR4), and erase status (SR5) bits in the status register to "1," but it cannot clear them to "0." Issuing this command clears those bits to "0."
60h	Protection Configuration Setup	First	Prepares the CSM for changes to the block locking status. If the next command is not BLOCK UNLOCK, BLOCK LOCK, or BLOCK LOCK DOWN, the command will be ignored, and the device will go to read status mode.
70h	Read Status Register	First	Places the device into read status register mode. Reading the device will output the contents of the status register for the addressed bank. The device will automatically enter this mode for the addressed bank after a PROGRAM or ERASE operation has been initiated.
90h	Read Protection Configuration	First	Puts the device into the read protection configuration mode so that reading the device will output the manufacturer/device codes or block lock status.
98h	Read Query	First	Puts the device into read query mode so that reading the device will output common Flash interface information.
B0h	Program/Erase Suspend	First	Suspends the currently executing PROGRAM/ERASE operation. The status register will indicate when the operation has been successfully suspended by setting either the program suspend (SR2) or erase suspend (SR6), and the WSM status bit (SR7) to a "1" (ready). The WSM will continue to idle in the suspend state, regardless of the state of all input control signals except RST#, which will immediately shut down the WSM and the remainder of the chip if RST# is driven to VIL.
C0h	Program Device Protection Register	First	Writes a specific code into the device protection register.
	Lock Device Protection Register	First	Locks the device protection register; data can no longer be changed.
D0h	Erase Confirm	Second	If the previous command was an ERASE SETUP command, then the CSM will close the address and data latches, and it will begin erasing the block indicated on the address bus. During programming/erase, the device will respond only to the READ STATUS REGISTER, PROGRAM SUSPEND, or ERASE SUSPEND commands and will output status register data on the falling edge of OE# or CE#, whichever occurs last.
	Program/Erase Resume	First	If a PROGRAM or ERASE operation was previously suspended, this command will resume the operation.
D0h	FPA Confirm	Second	If the previous command was FPA SETUP, the CSM will latch the address indicated on the address bus and enter the FPA mode.
FFh	Read Array	First	During the array mode, array data will be output on the data bus.



Table 5: Command Descri	ptions (Continued)
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CODE	DEVICE MODE	BUS CYCLE	DESCRIPTION
01h	Lock Block	Second	If the previous command was PROTECTION CONFIGURATION SETUP, the CSM will latch the address and lock the block indicated on the address bus.
2Fh	Lock Down	Second	If the previous command was PROTECTION CONFIGURATION SETUP, the CSM will latch the address and lock down the block indicated on the address bus.
D0h	Unlock Block	Second	If the previous command was PROTECTION CONFIGURATION SETUP, the CSM will latch the address and unlock the block indicated on the address bus. If the block had been previously set to lock down, this operation will have no effect.
00h	Invalid/Reserved		Unassigned command that should not be used.

Clear Status Register

The internal circuitry can set, but not clear, the block lock status bit (SR1), the VPP status bit (SR3), the program status bit (SR4), and the erase status bit (SR5) of the status register. The CLEAR STATUS REGISTER command (50h) allows the external microprocessor to clear these status bits and synchronize to the internal operations. When the status bits are cleared, the device returns to the read array mode.

Read Operations

The following READ operations are available: READ ARRAY, READ PROTECTION CONFIGURATION REG-ISTER, READ QUERY, and READ STATUS REGISTER.

Read Array

The array is read by entering the command code FFh on DQ0–DQ7. Control signals CE# and OE# must be at a logic LOW level (VIL), and WE# and RST# must be at logic HIGH level (VIH) to read data from the array. Data is available on DQ0–DQ15. Any valid address within any of the blocks selects that address and allows data to be read from that address. Upon initial power-up or device reset, the device defaults to the read array mode.

Read Protection Configuration Data

The chip identification mode outputs three types of information: the manufacturer/device identifier, the block locking status, and the protection register. Two bus cycles are required for this operation: the chip identification data is read by entering the command code 90h on DQ0–DQ7 to the bank containing address 0h and the identification code address on the address lines. Control signals CE# and OE# must be at a logic LOW level (VIL), and WE# and RST# must be at a logic HIGH level (VIH) to read data from the protection configuration register. Data is available on DQ0–DQ15. After data is read from the protection configuration register, the read array command, FFh, must be issued to the bank containing address 00h prior to performing other operations. See Table 10 for further details.

Read Query

The read query mode outputs common flash interface (CFI) data when the device is read (see Table 12). Two bus cycles are required for this operation. It is possible to access the query by writing the read query command code 98h on DQ0–DQ7 to the bank containing address 0h. Control signals CE# and OE# must be at a logic LOW level (VIL), and WE# and RST# must be at a logic HIGH level (VIH) to read data from the query. The CFI data structure contains information such as block size, density, command set, and electrical specifications. After reading CFI data, the READ ARRAY command (FFh) must be issued to the bank containing address 00h prior to performing other operations.

Read Status Register

The status register is read by entering the command code 70h on DQ0–DQ7. Two bus cycles are required for this operation: one to enter the command code and the block address and a second to read the status register. In a READ cycle, the address is latched and register data is updated on the falling edge of OE# or CE#, whichever occurs last.



Table 6: Command State Machine Transition Table

	nt	of Prese		Prese Parti			n)	t partitic	the presen	ext state of	n (and ne	nt partitio	the prese	input to	Command	С									
Present State of the other Partition	Mode	State	Data When Read	SR7	FFh Read Array	30h FPA Setup	10h/40h Program Setup	20h Erase Setup	D0h BE Confirm, P/E Resume, ULB Confirm	B0h Program /Erase Suspend	70h Read Status	50h Clear Status Register	90h Read Device ID	98h Read Query	60h Lock/ Unlock/ Lock Down	C0h OTP Setup	01h Lock Confirm	2Fh Lock Down Confirm							
1 Setup 2 Busy							array	Read									ead array	R							
3 Idle								Erase setup	array	Read	Read	Read	Read	Read		OTP setup	array	Read							
4 Erase Susp		Array	Array		ID	query		setup																	
5 Prog							array	Read									ead array	R							
5 Susp 6 Setur	-							Read									ead array	R							
7 Busy							unuy	Erase								OTP									
8 Idle		Query	CFI	1	Read array	n setup	Progra	setup	array	Read	Read status	Read array	Read ID	Read query		setup	array	Read							
9 Erase Susp					Read array									ead array	R										
10 Prog Susp					Read array																				
11 Setup	Read														Lock/										
12 Busy		Device			Read array				Read	Read	Read	Read	Unlock/ Lock down		ead array	R									
13 Idle		ID	ID	1	Read		_	Erase setup	array	Read	status	array								ID	query		OTP setup	array	Read
14 Erase Susp				Read array array	Program setup array																				
15 Prog Susp							array	Read									ead array	R							
16 Setup							arrav	Read									ead array	R							
17 Busy 18 Idle								Erase	array	Bood						OTP	array								
Eroco		Status	Status	1	Read array	n setup	Progra	setup			Read status	Read array	Read ID	Read query		setup	array	Redu							
19 Susp									Read array								ead array	R							
20 Susp		Catura	Status	1			array	Read				Ductosticu													
21 Idle 22 Idle	-	Setup Busy	Status	1							-	Protection Protection													
23 Setup							array	Read		-	-				Lock/ Unlock/ Lock down		ead array	R							
24 Busy	Prot reg	D.	Ch. i		<u> </u>			F			Read	Read	Read	Read		075									
25 Idle		Done	Status	1	Read	n setup	Progra	Erase setup	array	Read	status	array	ID	query		OTP setup	array	Read							
26 Erase Susp					array		, , , , , , , , , , , , , , , , , , ,		Read array	F							ead array	R							
27 Prog Susp							array	Read									,								



Table 6: Command State Machine Transition Table (Continued)

		C	Command	input to	the prese	nt partitio	n (and ne	ext state of	the presen	t partitic	n)			Prese Parti		of Preser	nt		
Fh ock own onfirm	01h Lock Confirm	C0h OTP Setup	60h Lock/ Unlock/ Lock Down	98h Read Query	90h Read Device ID	50h Clear Status Register	70h Read Status	B0h Program /Erase Suspend	D0h BE Confirm, P/E Resume, ULB Confirm	20h Erase Setup	10h/40h Program Setup	30h FPA Setup	FFh Read Array	SR7	Data When Read	State	Mode		nt State e other ion
LB/L	ULB				Read Sta	tus			LB/ULB		Read	status		1	Status	Setup		28	Any state
Re	ead array	•								Read	array							29 30	Setup Busy
Read	array	OTP		Bood	Road	Pood	Pood		Read array		İ							31	Idle
		setup	-	Read query	Read ID	Read array	Read status		rase setup Read array		Progra	m setup	Read array	1	Status	Error		32	Erase
R	ead array		Lock/							Read	array						Unlock/ lock	33	susp Prog
			Unlock/ Lock								-						down	34	susp Setup
Re	ead array		down							Read	array							35	Busy
Read	array	OTP setup		Read	Read	Read	Read		Read array Trase setup		Progra	m setup	Read	1	Status	Lock/		36	Idle
D				query	ID	array	status	I	Read array		riogra	in setup	array			Unlock		37	Erase susp
ĸ	ead array									Read	array							38	Prog susp
						Progr	am Busy						1	Status	Setup		39	Any state	
			Program	n Busy				PS read status		Р	rogram bus	5y		0	Status	Busy		40	Idle
R	ead array									Read	arrav				Read	array		41	Setup
	,	ΟΤΡ							Read array		,		1			,	Prog	42	Busy
Read	array	setup	-	Read query	Read ID	Read array	Read status		rase setup		Progra	m setup	Read array	1	Read	l array		43	Idle
Re	ead array							I	Read array				array		Read	l array		44	Erase susp
										Read	array					-		45	Prog susp
				Prog	Prog	Prog	Prog	Prog										46	Setup
Progra	am Read a	rray		susp read	susp read	susp read	susp read	susp read	Program busy	Pro	gram suspe	end read a	rray	1	Status	Read status		47	Idle Erase
			Lock/ Unlock/	query	ID	array	status	array										48	susp
_			Lock down	Prog	Prog	Prog	Prog	Prog										49	Setup
Program	m suspend array	read		susp read query	susp read ID	susp read array	susp read status	susp read array	Program busy	Pro	gram suspe	end read ai	rray	1	Array	Read array		50 51	Idle Erase
						-		-									Prog susp	52	susp Setup
Program	m suspend	read		Prog susp	Prog susp	Prog susp	Prog susp	Prog susp	Program	D.: -		and read		1	ID	Read		53	Idle
-	array			read query	read ID	read array	read status	read array	busy	Pro	gram suspe	end read al	ıdy		U	ID		54	Erase susp
				Prog	Prog	Prog	Prog	Prog										55	Setup
Program	m suspend	read		susp	susp	susp	susp	susp	Program	Pro	gram suspe	end read a	rray	1	CFI	Read		56	Idle
	array			read query	read ID	read array	read status	read array	busy		- '					query		57	Erase susp



Table 6: Command State Machine Transition Table (Continued)

		C	Command	input to	the prese	nt partitio	n (and ne	ext state of	the presen	t partitic	n)			Pres Part		of Prese	nt		
2Fh Lock Down Confirm	01h Lock Confirm	C0h OTP Setup	60h Lock/ Unlock/ Lock Down	98h Read Query	90h Read Device ID	50h Clear Status Register	70h Read Status	B0h Program /Erase Suspend	D0h BE Confirm, P/E Resume, ULB Confirm	20h Erase Setup	10h/40h Program Setup	30h FPA Setup	FFh Read Array	SR7	Data When Read	State	Mode		nt State e other ion
			F	Read arra	iy				Erase busy		Read	array		1	Status	Setup		58	Idle
F	Read array									Read	array							59	Setu
Read	d array	OTP							Read array									60 61	Busy Idle
neue	andy	setup		Read query	Read ID	Read array	Read status		Frase setup Read array		Progra	m setup	Read array	1	Status	Error			Eras
F	Read array		Lock/						kead array									62	susp Proc
			Unlock/ Lock							Read	array						Erase	63 64	susp
F	Read array		down							Read	array							65	Setu Busy
Read	d array	OTP setup		Read	Read	Read	Read		Read array Frase setup				Read	1	Status	Dono		66	Idle
				query	ID	array	status	1	Read array		Prograi	m setup	array		Status	Done		67	Erase
F	Read array									Read	array							68	Prog susp
			Block era	ise busy				ES read status			Erase busy			0	Status	Busy		69	Idle
								ES read	Erase	E	rase suspen	d read arr	av					70	Setu
				Erase	Erase	Erase	Erase	array	busy Eras		d read arra		.,			Read status		71	Busy
Erase su	uspend read	l array		susp read query	susp read ID	susp read array	susp read status	ES read array	Erase busy	ES read array	Progra	m setup	ES read array	1	Status			72	Idle
									Eras	se suspen	d read arra	iy						73	Prog
								ES read array	Erase busy	E	rase suspen	d read arr	ау					74	Setu
				Erase	Erase	Erase	Erase	anay		se suspen	d read arra	у						75	Busy
Erase su	uspend read	l array		susp read query	susp read ID	susp read array	susp read status	ES read array	Erase busy	ES read array	Progra	m setup	ES read array	1	Array	Read array		76	Idle
			Lock/ Unlock/						Eras	se suspen	d read arra	у					Erase	77	Prog
			Lock down					ES read array	Erase busy	E	rase suspen	d read arr	ау				Susp	78	Setu
				Erase	Erase	Erase	Erase			se suspen	d read arra	iy						79	Busy
Erase su	uspend read	l array		susp read query	susp read ID	susp read array	susp read status	ES read array	Erase busy	ES read array	Progra	m setup	ES read array	1	ID	Read ID		80	Idle
									Eras	se suspen	d read arra	iy]				81	Prog
								ES read array	Erase busy	E	rase suspen	d read arr	ау					82	Setu
				Erase	Erase	Erase	Erase			se suspen	d read arra	iy						83	Busy
Erase su	uspend read	l array		susp read query	susp read ID	susp read array	susp read status	ES read array	Erase busy	ES read array	Progra	m setup	ES read array	1	CFI	Read query		84	Idle
									Eras	se suspen	d read arra	iy						85	Prog susp



Programming Operations

A PROGRAM operation is iniated by a command sequence to the CSM: PROGRAM SETUP (40h) or ALTERNATE PROGRAM SETUP (10h), followed by the correct address and data. Once the program command sequence has been entered, the CSM automatically enters into the status read mode. To access the status of the entered command, a READ of the device at the address given in the first cycle will provide the relevant status of the command. See Figure 4 on page 19 for details.

After the desired command code is entered (10h or 40h command code on DQ0–DQ7), the WSM takes over and correctly sequences the device to complete the PROGRAM operation. The WRITE operation may be monitored through the status register (see the Status Register section). During this time, the CSM will only respond to a PROGRAM SUSPEND command until the PROGRAM operation has been completed, after which time all commands to the CSM become valid again. After the PROGRAM operation is complete, the READ ARRAY command (FFh) must be given to the same partition to clear the CSM.

The PROGRAM operation can be suspended by issuing a PROGRAM SUSPEND command (B0h). Once the WSM reaches the suspend state, it allows the CSM to respond only to READ ARRAY, READ STATUS REG-ISTER, READ PROTECTION CONFIGURATION, READ QUERY, or PROGRAM RESUME. During the PROGRAM SUSPEND operation, array data should be read from an address other than the one being programmed. To resume the PROGRAM operation, a PROGRAM RESUME command (D0h) must be issued to the suspended partition, causing the CSM to clear the suspend state previously set (see Figure 4 for programming operation and Figure 5 for program suspend and program resume). The RESUME command (D0h) must not be given to a ready partition when the other partition is suspended. Doing so could cause the suspended partition to resume.

Taking RST# to VIL during programming aborts the PROGRAM operation. During programming, VPP must remain in the appropriate VPP voltage range as shown in the recommended operating conditions table.

Fast Programming Algorithm (FPA) Mode

The fast programming algorithm (FPA) is intended for in-factory use. It enables fast data stream programming.

2 MEG x 16 ASYNC/PAGE FLASH MEMORY

For in-factory programming, the FPA, along with an optimized set of programming parameters, minimizes chip programming time when 11.4V < VPP 12.6V.

Executing the FPA command (30h), followed by FPA CONFIRM (D0h), enables an entire block to be programmed. This eliminates the need to continuously update the address to be programmed.

An initial delay is required after issuing the FPA command. (See the Erase and Program Cycle Timing Requirements Table.) The delay enables the device to detect 12V on VPP. If VPP < 11.4V, or if the block is locked, the status register returns an error. When the FPA command is executed successfully, a data stream can be programmed beginning at the first address. The address can be held constant, or it can be incremented within the address range. The program ends when the programmer enters an address outside the address range of the current block.

When the FPA is activated, the data must be provided in sequential order to the WSM. Immediately after programming, verification is executed. The address sequence is again provided to the WSM, which automatically performs a data consistency check between the data stored in the memory array, and the programmed data. The result is stored in the status register. Issuing an address outside the memory block boundaries exits the verification cycle. Figure 8 shows the FPA flowchart.

Erase Operations

An ERASE operation must be used to initialize all bits in an array block to "1s." After BLOCK ERASE confirm is issued, the CSM responds only to an ERASE SUSPEND command until the WSM completes its task.

Block erasure inside the memory array sets all bits within the address block to logic 1s. Erase is accomplished only by blocks; data at single address locations within the array cannot be erased individually. The block to be erased is selected by using any valid address within that block. Block erasure is initiated by a command sequence to the CSM: BLOCK ERASE setup (20h) followed by BLOCK ERASE CONFIRM (D0h) (see Figure 6). A two-command erase sequence protects against accidental erasure of memory contents.

Once the BLOCK ERASE CONFIRM command is entered, the CSM automatically enters into status read mode. To access the status of the entered command, a READ of the device at the address given in the first cycle will provide the relevant status of the command. See Figure 6 on page 21 for details.



When the BLOCK ERASE CONFIRM command is complete, the WSM automatically executes a sequence of events to complete the block erasure. During this sequence, the block is programmed with logic 0s, data is verified, all bits in the block are erased, and finally verification is performed to ensure that all bits are correctly erased. The ERASE operation may be monitored through the status register (see the Status Register section). After the ERASE operation is complete, the READ ARRAY command (FFh) must be given to the same partition to clear the CSM.

During the execution of an ERASE operation the ERASE SUSPEND command (B0h) can be entered to direct the WSM to suspend the ERASE operation. Once the WSM has reached the suspend state, it allows the CSM to respond only to the READ ARRAY, READ STA-TUS REGISTER, READ QUERY, READ CHIP PROTEC- TION CONFIGURATION, PROGRAM SETUP, PROGRAM RESUME, ERASE RESUME and LOCK SETUP (see the Block Locking section). During the ERASE SUSPEND operation, array data must be read from a block other than the one being erased. To resume the ERASE operation, an ERASE RESUME command (D0h) must be issued to cause the CSM to clear the suspend state previously set (see Figure 7). It is also possible to suspend an ERASE in any bank and initiate a WRITE to another block in the same bank. After the completion of a WRITE, an ERASE can be resumed by writing an ERASE RESUME command. The RESUME command (D0h) must not be given to a ready partition when the other partition is suspended. Doing so could cause the suspended partition to resume.

lable /:	Bus Operations	

MODE	RST#	CE#	OE#	WE#	ADDRESS	DQ0-DQ15
Read (array, status registers, device identification register, or query)	Viн	VIL	VIL	Viн	X	Dout
Standby	Vih	Vih	Х	Х	Х	High-Z
Output Disable	VIH	Vih	Х	Х	Х	High-Z
Reset	VIL	Х	Х	Х	Х	High-Z
Write	VIL	VIL	Viн	VIL	Х	DIN



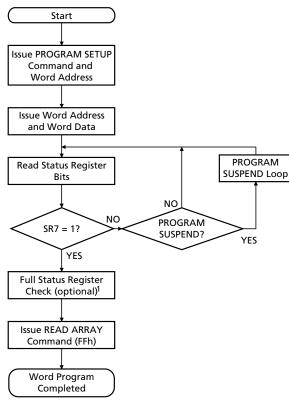
Table 8: Status Register Bit Definition

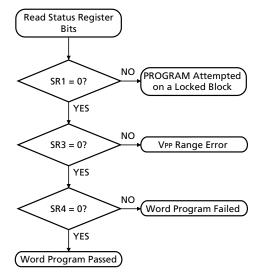
WSMS	ESS	ES	PS	VppS	PSS	BLS	FPAS
7	6	5	4	3	2	1	0

STATUS BIT #	STATUS REGISTER BIT	DESCRIPTION
SR7	WRITE STATE MACHINE STATUS (WSMS) 1 = Ready 0 = Busy	Check write state machine bit first to determine word program or block erase completion, before checking program or erase status bits.
SR6	ERASE SUSPEND STATUS (ESS) 1 = BLOCK ERASE Suspended 0 = BLOCK ERASE in Progress/Completed	When ERASE SUSPEND is issued, WSM halts execution and sets both WSMS and ESS bits to "1." ESS bit remains set to "1" until an ERASE RESUME command is issued.
SR5	ERASE STATUS (ES) 1 = Error in Block Erasure 0 = Successful BLOCK ERASE	When this bit is set to "1," WSM has applied the maximum number of erase pulses to the block and is still unable to verify successful block erasure.
SR4	PROGRAM STATUS (PS) 1 = Error in PROGRAM 0 = Successful PROGRAM	When this bit is set to "1," WSM has attempted but failed to program a word.
SR3	VPP STATUS (VPPS) 1 = VPP Low Detect, Operation Abort 0 = VPP = OK	The VPP status bit does not provide continuous indication of the VPP level. The WSM interrogates the VPP level only after the program or erase command sequences have been entered and informs the system if VPP < 1.8V. The VPP level is also checked before the PROGRAM/ERASE operation is verified by the WSM.
SR2	PROGRAM SUSPEND STATUS (PSS) 1 = PROGRAM Suspended 0 = PROGRAM in Progress/Completed	When PROGRAM SUSPEND is issued, WSM halts execution and sets both WSMS and PSS bits to "1." PSS bit remains set to "1" until a PROGRAM RESUME command is issued.
SR1	BLOCK LOCK STATUS (BLS) 1 = PROGRAM/ERASE Attempted on a Locked Block; Operation Aborted 0 = No Operation to Locked Blocks	If a PROGRAM or ERASE operation is attempted to one of the locked blocks, this is set by the WSM. The operation specified is aborted and the device is returned to read status mode.
SRO	FAST PROGRAMMING ALGORITHM STATUS (FPAS) 1 = FPA PROGRAM/ERASE Busy 0 = FPA Ready	When this bit is set to "1," the FPA algorithm is active. When the FPA operation is complete, this bit is reset to "0."



Figure 4: Automated Word Programming Flowchart





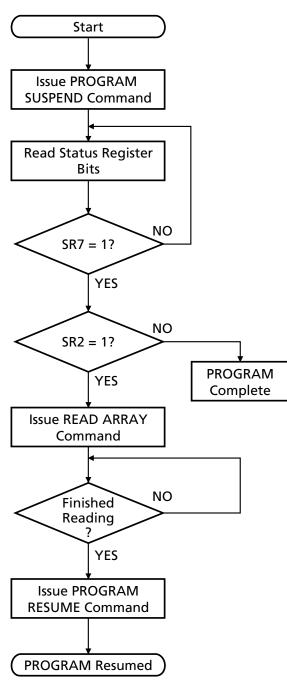
BUS OPERATION	COMMAND	COMMENTS				
WRITE	WRITE	Data = 40h or 10h				
	PROGRAM	Addr = Address of word to				
	SETUP	be programmed				
WRITE	WRITE	Data = Word to be				
	DATA	programmed				
		Addr = Address of word to				
		be programmed				
READ		Status register data				
		Toggle OE# or CE# to				
		update status register.				
Standby		Check SR7				
		1 = Read, 0 = Busy				
WRITE	READ	Data = FFh				
	ARRAY	Addr = Bank of word to				
		be programmed				
Repeat for subsequent words.						
Write FFh after each word programming operation to						
clear the CSM a	and reset the d	levice to read array mode.				

BUS OPERATION	COMMAND	COMMENTS
Standby		Check SR1 1 = Detect locked block
Standby		Check SR3 ² 1 = Detect VPP LOW
Standby		Check SR4 ³ 1 = Word program error

- NOTE: 1. Full status register check can be done after each word or after a sequence of words.
 - 2. SR3 must be cleared before attempting additional PROGRAM/ERASE operations.
 - 3. SR4 is cleared only by the CLEAR STATUS REGISTER command, but it does not prevent additional program operation attempts.



Figure 5: PROGRAM/SUSPEND/ PROGRAM/RESUME Flowchart

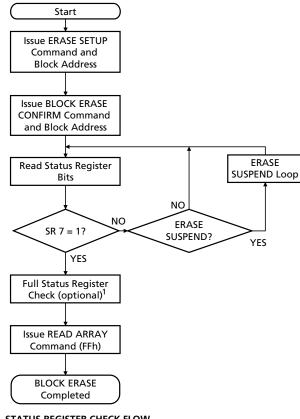


BUS OPERATION	COMMAND	COMMENTS
WRITE	PROGRAM SUSPEND	Data = B0h
READ		Status register data Toggle OE# or CE# to update status register
Standby		Check SR7 1 = Ready
Standby		Check SR2 1 = Suspended
WRITE	READ ARRAY	Data = FFh
READ		Read data from block other than that being programmed.
WRITE	PROGRAM RESUME	Data = D0h

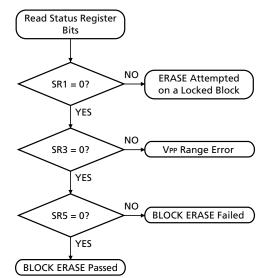
- NOTE: 1. Full status register check can be done after each word or after a sequence of words.
 - 2. SR3 must be cleared before attempting additional PROGRAM/ERASE operations.
 - 3. SR5 is cleared only by the CLEAR STATUS REGISTER command in cases where multiple blocks are erased before full status is checked.



Figure 6: BLOCK ERASE Flowchart



FULL STATUS R	EGISTER CHECK FLOW
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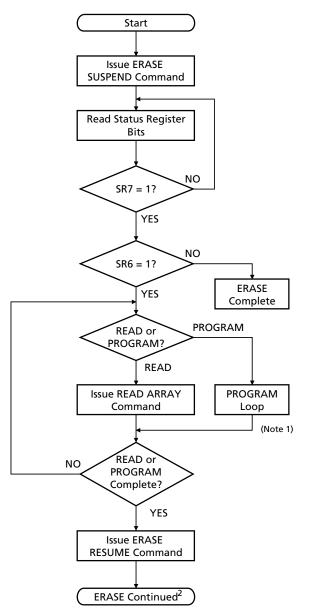
BUS OPERATION	COMMAND	COMMENTS		
WRITE	WRITE ERASE SETUP	Data = 20h Block Addr = Address within block to be erased		
WRITE	ERASE	Data = D0h Block Addr = Address within block to be erased		
READ		Status register data Toggle OE# or CE# to update status register.		
Standby		Check SR7 1 = Read, 0 = Busy		
WRITE	READ ARRAY	Data = FFh Block Addr = Bank of word to be programmed		
Repeat for subsequent words Write FFh after each BLOCK ERASE operation to clear the CSM and reset the device to read array mode.				

BUS OPERATION	COMMAND	COMMENTS
Standby		Check SR1 1 = Detect locked block
Standby		Check SR3 ² 1 = Detect VPP block
Standby		Check SR5 ³ 1 = BLOCK ERASE error

- NOTE: 1. Full status register check can be done after each block or after a sequence of blocks.
 - 2. SR3 must be cleared before attempting additional PROGRAM/ERASE operations.
 - 3. SR5 is cleared only by the CLEAR STATUS REGISTER command in cases where multiple blocks are erased before full status is checked.



Figure 7: ERASE SUSPEND/ERASE RESUME Flowchart



BUS OPERATION	COMMAND	COMMENTS
WRITE	ERASE SUSPEND	Data = B0h
READ		Status register data Toggle OE# or CE# to update status register
Standby		Check SR7 1 = Ready
Standby		Check SR6 1 = Suspended
WRITE	READ ARRAY	Data = FFh
READ		Read data from block other than that being erased.
WRITE	ERASE RESUME	Data = D0h

NOTE: 1. See BLOCK ERASE Flowchart for complete erasure procedure.

2. See Word Programming Flowchart for complete programming procedure.



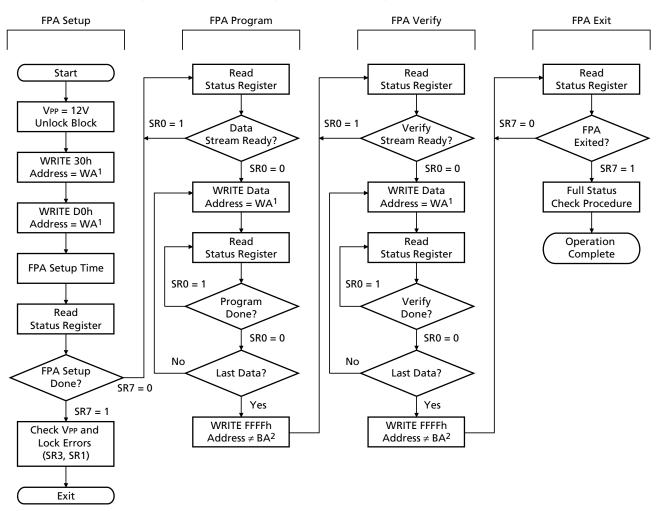


Figure 8: Fast Programming Algorithm Flowchart

Note: 1. WA: Word address of memory location to be written.

First address may be held constant or incremented within the address range.

2. BA: Address within the block.



Read-While-Write/Erase Concurrency

It is possible for the device to read from one bank while erasing/writing to another bank. Once a bank enters the WRITE/ERASE operation, the other bank automatically enters read array mode. For example, during a READ CONCURRENCY operation, if a PRO-GRAM/ERASE command is issued in bank a, then bank *a* changes to the read status mode and bank *b* defaults to the read array mode. The device will read from bank b if the latched address resides in bank b (see Figure 9). Similarly, if a PROGRAM/ERASE command is issued in bank b, then bank b changes to read status mode and bank *a* defaults to read array mode. When returning to bank *a*, the device will read PRO-GRAM/ERASE status if the latched address resides in bank a. A correct bank address must be specified to read status register after returning from concurrent read in the other bank.

When reading the CFI or the chip protection register, concurrent operation is not allowed on the top boot device. Concurrent READ of the CFI or the chip protection register is only allowed when a PROGRAM or ERASE operation is performed on bank b on the bottom boot device. For a top boot device, reading of the CFI table or the chip protection register is only allowed if bank *a* is in read array mode.

Figure 9: Read-While-Write Concurrency

Bank a	Bank b
 Erasing/writing to bank a Erasing in bank a can be suspended, and a WRITE to another block in bank a can be initiated. After the WRITE in that block is complete, an ERASE can be resumed by writing an ERASE RESUME command. 	1 - Reading from bank <i>b</i>
1 - Reading bank a	 Erasing/writing to bank b Erasing in bank b can be suspended, and a WRITE to another block in bank b can be initiated. After the WRITE in that block is complete, an ERASE can be resumed by writing an ERASE RESUME command.



Block Locking

The MT28F322P3 Flash memory provides a flexible locking scheme which allows each block to be individually locked or unlocked with no latency.

The device offers two-level protection for the blocks. The first level allows software-only control of block locking (for data which needs to be changed frequently), while the second level requires hardware interaction before locking can be changed (code which does not require frequent updates).

Control signals WP#, DQ0, and DQ1 define the state of a block; for example, state [001] means WP# = 0, DQ0 = 0 and DQ1 = 1.

Table 9 defines all of the possible locking states.

Locked State

After a power-up sequence completion, or after a reset sequence, all blocks are locked (states [001] or [101]). This means full protection from alteration. Any PROGRAM or ERASE operations attempted on a locked block will return an error on bit SR1 of the status register. The status of a locked block can be changed to unlocked or lock down using the appropriate software commands. Writing the lock command sequence, 60h followed by 01h, can lock an unlocked block.

Unlocked State

Unlocked blocks (states [000], [100], [110]) can be programmed or erased. All unlocked blocks return to the locked state when the device is reset or powered down. An unlocked block can be locked or locked down using the appropriate software command sequence, 60h followed by 01h or 2Fh, respectively. (see Table 4). Locked down blocks revert to the locked state when the device is reset or powered down (see Table 4).

Locked Down State

Blocks that are locked down (state [011]) are protected from PROGRAM and ERASE operations, but their protection status cannot be changed using software commands alone. A locked or unlocked block can be locked down by writing the lock down command sequence, 60h followed by 2Fh. Locked down blocks revert to the locked state when the device is reset or powered down.

The LOCK DOWN function is dependent on the WP# input. When WP# = 0, blocks in lock down [011] are protected from program, erase, and lock status changes. When WP# = 1, the lock down function is disabled ([111]), and locked down blocks can be individually unlocked by a software command to the [110] state, where they can be erased and programmed. These blocks can then be relocked [111] and unlocked [110] as desired while WP# remains HIGH. When WP# goes LOW, blocks that were previously locked down return to the locked down state [011] regardless of any changes made while WP# was HIGH. Device reset or power-down resets all locks, including those in lock down, to locked state (see Table 10).

Reading a Block's Lock Status

The lock status of every block can be read in the read device identification mode. To enter this mode, write 90h to the bank containing address 00h. Subsequent READs at block address +00002 will output the lock status of that block. The lowest two outputs, DQ0 and DQ1, represent the lock status. DQ0 indicates the block lock/unlock status and is set by the LOCK com-

WP#	DQ1	DQ0	NAME	ERASE/PROGRAM ALLOWED	LOCK	UNLOCK	LOCK DOWN
0	0	0	Unlocked	Yes	To [001]	No Change	To [011]
0	0	1	Locked (Default)	No	No Change	To [000]	To [011]
0	1	1	Lock Down	No	No Change	No Change	No Change
1	0	0	Unlocked	Yes	To [101]	No Change	To [111]
1	0	1	Locked	No	No Change	To [100]	To [111]
1	1	0	Lock Down Disabled	Yes	To [111]	No Change	To [111]
1	1	1	Lock Down Disabled	No	No Change	To [110]	No Change

Table 9: Block Locking State Transition

NOTE: All blocks are software-locked upon power-up sequence completion.



mand and cleared by the UNLOCK command. It is also automatically set when entering lock down. DQ1 indicates lock down status and is set by the LOCK DOWN command. It can only be cleared by reset or powerdown, not by software. Table 9 shows the locking state transition scheme. The read array command, FFh, must be issued to the bank containing address 00h prior to performing other operations.

Locking Operations During Erase Suspend

Changes to block lock status can be performed during an ERASE SUSPEND by using the standard locking command sequences to unlock, lock, or lock down. This is useful in the case when another block needs to be updated while an ERASE operation is in progress.

To change block locking during an ERASE operation, first write the ERASE SUSPEND command (B0h), then check the status register until it indicates that the ERASE operation has been suspended. Next, write the desired lock command sequence to block lock, and the lock status will be changed. After completing any desired LOCK, READ, or PROGRAM operations, resume the ERASE operation with the ERASE RESUME command (D0h).

If a block is locked or locked down during an ERASE SUSPEND on the same block, the locking status bits will be changed immediately. When the ERASE is resumed, the ERASE operation will complete.

A locking operation cannot be performed during a PROGRAM SUSPEND.

Chip Protection Register

A 128-bit chip protection register can be used to fulfill the security considerations in the system (preventing the device substitution).

The 128-bit security area is divided into two 64-bit segments. The first 64 bits are programmed at the manufacturing site with a unique 64-bit unchangeable number. The other segment is left blank for customers to program as desired. (See Figure 10).

Reading the Chip Protection Register

The chip protection register is read in the device identification mode. To enter this mode, load the 90h command to the bank containing address 00h. Once in this mode, READ cycles from addresses shown in Table 10 retrieve the specified information. To return to the read array mode, write the READ ARRAY command (FFh). The READ ARRAY command, FFh, must be issued to the bank containing address 00h prior to performing other operations.

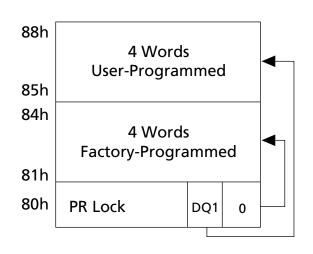
Programming the Chip Protection Register

The first 64 bits (PR1) of the protection register (addresses 81h–84h) are programmed with a unique identifier at the factory. DQ0 of the PR lock register (address 80h) is programmed to a "0" state, locking the first 64 bits and preventing any further programming.

The second 64 bits (PR2) are a user area (addresses 85h–88h), where the user can program any information into this area as long as DQ1 of the PR lock register remains unprogrammed. After DQ1 of the PR lock register is programmed, no further programming is allowed on PR2. The programming sequence is similar to array programming except that the PROTECTION REGISTER PROGRAMMING SETUP command (C0h) is issued instead of an ARRAY PROGRAMMING SETUP command (40h), followed by the data to be programmed at addresses 85h–88h.

To program the PR lock bit for PR2 (to prevent further programming), use the above sequence on address 80h, with data FFFDh (DQ1 = 0).

Figure 10: Protection Register Memory Map





ITEM	ADDRESS ²	DATA
Manufacturer Code (x16)	00000h	002Ch
Device CodeTop boot configurationBottom boot configuration	00001h	4494h 4495h
Block Lock ConfigurationBlock is unlockedBlock is lockedBlock is locked down	XX002h	Lock DQ0 = 0 DQ0 = 1 DQ0 = 1
Chip Protection Register Lock	80h	PR Lock
Chip Protection Register 1	81h–84h	Factory Data
Chip Protection Register 2	85h–88h	User Data

 Table 10:
 Chip Configuration Addressing¹

NOTE:1.Other locations within the configuration address space are reserved by Micron for future use. 2."XX" specifies the block address of lock configuration.

Asynchronous Read Cycle

When accessing addresses in a random order or when switching between pages, the access time is given by ^tAA.

When CE# and OE# are LOW, the data is placed on the data bus and the processor can read the data.

Page Read Mode

The initial portion of the page mode cycle is the same as the asynchronous access cycle. Holding CE# LOW and toggling addresses A0–A2 allows random access of other words in the page. The page word size is eight words.

VPP/Vcc Program and Erase Voltages

The MT28F322P3 Flash memory provides in-system programming and erase with VPP in the 1.8V–3.3V range. The 12V VPP mode programming is offered for compatibility with existing programming equipment, but does not enhance programming performance using the standard programming commands.

The device can withstand 100,000 WRITE/ERASE operations when VPP = VCC or 100 WRITE/ERASE operations and 10 cumulative hours when VPP = 12V.

In addition to the flexible block locking, the VPP programming voltage can be held LOW for absolute hardware write protection of all blocks in the Flash device. When VPP is below VPPLK, any PROGRAM or ERASE operation will result in an error, prompting the corresponding status register bit (SR3) to be set. During WRITE and ERASE operations, the WSM monitors the VPP voltage level. WRITE/ERASE operations are allowed only when VPP is within the ranges specified in Table 11.

When VCC is below VLKO, any WRITE/ERASE operation will be disabled.

Table 11: VPP Range (V)

	MIN	MAX
In-System	1.8	3.3
In-Factory	11.4	12.6

Standby Mode

ICC supply current is reduced by applying a logic HIGH level on CE# and RST# to enter the standby mode. In the standby mode, the outputs are placed in High-Z. Applying a CMOS logic HIGH level on CE# and RST# reduces the current to ICC3 (MAX). If the device is deselected during an ERASE operation or during programming, the device continues to draw current until the operation is complete.

Automatic Power Save (APS) Mode

Substantial power savings are realized during periods when the array is not being read and the device is in the active mode. During this time the device switches to the automatic power save mode. When the device switches to this mode, ICC is reduced to a level comparable to ICC3. Further power savings can be real-



ized by applying a logic HIGH level on CE# to place the device in standby mode. The low level of power is maintained until another operation is initiated. In this mode, the I/Os retain the data from the last memory address read until a new address is read. This mode is entered automatically if no address or control signals toggle.

Device Reset

To correctly reset the MT28F322P3 Flash memory, the RST# signal must be asserted (RST# = VIL) for a minimum of ^tRP. After reset, the device can be accessed for a READ operation with a delayed access time of ^tRWH from the rising edge of RST#. The circuitry used for generating the RST# signal needs to be

common with the rest of the system reset to ensure that correct system initialization occurs. Please refer to the timing diagram for further details.

Power-Up Sequence

The following power-up sequence is recommended to initialize internal chip operations:

- At power-up, RST# should be kept at VIL for 2µs after VCC reaches VCC (MIN).
- VCCQ should not come up before VCC.
- VPP should be kept at VIL to maximize data integrity.

When the power-up sequence is completed, RST# should be brought to VIH. To ensure a proper power-up, the rise time of RST# (10%–90%) should be <10µs.



Absolute Maximum Ratings*

Voltage to Any Ball Except VCC and VPI	0
with Respect to Vss	0.5V to +4V
VPP Voltage (for BLOCK ERASE and PR	OGRAM
with Respect to Vss	-0.5V to +13.5V**
VCC and VCCQ Supply Voltage	
with Respect to Vss	0.3V to +4V
Output Short Circuit Current	100mA
Operating Temperature Range	40°C to +85°C
Storage Temperature Range	55°C to +150°C
Operating Temperature Range	40°C to +85°C
Soldering Cycle	260°C for 10s

2 MEG x 16 ASYNC/PAGE FLASH MEMORY

*Stresses greater than those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.

**Maximum DC voltage on VPP may overshoot to +13.5V for periods <20ns.

PARAMETER SYMBOL MIN MAX UNITS NOTES Operating temperature -40 +85 °C tΑ Vcc supply voltage Vcc 2.7 3.3 V 3.3 V I/O supply voltage VccO 2.2 VPP voltage VPP1 1.8 3.3 V VPP in-factory programming voltage Vpp2 11.4 12.6 V 100,000 Block erase cycling VPP = VPP1VPP1 _ Cycles 100 VPP = VPP2VPP2 Cycles 1 _

Recommended Operating Conditions

NOTE: 1. VPP = VPP2 is a maximum of 10 cumulative hours.



DC Characteristics¹

PARAMETER	SYMBOL	MIN	ТҮР	MAX	UNITS	NOTES
Input Low Voltage	VIL	0	_	0.4	V	2
Input High Voltage	Vih	VccQ - 0.4	-	VccQ	V	2
Output Low Voltage IoL = 100μA	Vol	-0.10	_	0.10	V	
Output High Voltage Іон = -100µА	Voн	VccQ - 0.1	-	-	V	
VPP Lockout Voltage	VPPLK	-	-	1	V	
VPP During PROGRAM/ERASE	VPP1	1.8	-	3.3	V	
Operations	Vpp2	11.4	-	12.6	V	
Vcc Program/Erase Lock Voltage	Vlko	1	-	-	V	
Input Leakage Current	IL I	-	-	1	μA	
Output Leakage Current	loz	0.2	-	1	μΑ	
Vcc Read Current	Icc1					
Asynchronous Random Read, 70ns cycle Asynchronous Random Read, 200ns cycle			-	15 6	mA mA	3,4
Vcc Page Mode Read Current at 70ns/30ns	lcc2	-	_	7	mA	3,4
Vcc Standby Current	lcc3	-	15	50	μΑ	
Vcc Program Current	lcc4	_	10	20	mA	
Vcc Erase Current	Icc5	_	15	25	mA	
Vcc Erase Suspend Current	Icc6	_	15	50	μΑ	5
Vcc Program Suspend Current	lcc7	_	15	50	μΑ	5
Read-While-Write Current	lcc8	_	_	40	mA	
VPP Current (Read, Standby, Erase Suspend, Program Suspend)	IPP1					
VPP = VPP1 VPP = VPP2			-	1 200	μΑ μΑ	

NOTE: 1. All currents are in RMS unless otherwise noted.

2. VIL may decrease to -0.4V, and VIH may increase to VCCQ + 0.3V for durations not to exceed 20ns.

3. APS mode reduces Icc to approximately Icc3 levels.

4. Test conditions: Vcc = Vcc (MAX), CE# = VIL, OE# = VIH. All other inputs = VIH or VIL.

5. Icc6 and Icc7 values are valid when the device is deselected. Any READ operation performed while in suspend mode will have an additional current draw of suspend current (Icc6 or Icc7).



Capacitance

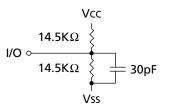
 $(^{T}A = +25^{\circ}C; f = 1 \text{ MHz})$

PARAMETER/CONDITION	SYMBOL	ТҮР	MAX	UNITS
Input Capacitance	C	7	12	pF
Output Capacitance	Cout	9	12	pF

Read Cycle Timing Requirements

		-70		-80		
PARAMETER	SYMBOL	MIN	MAX	MIN	MAX	UNITS
Address to output delay	^t AA		70		80	ns
CE# LOW to output delay	^t ACE		70		80	ns
Page address access	^t APA		30		30	ns
OE# LOW to output delay	^t AOE		25		25	ns
RST# HIGH to output delay	^t RWH		200		200	ns
RST# LOW pulse width	^t RP	100		100		ns
CE# or OE# HIGH to output High-Z	^t OD		25		25	ns
Output hold from address, CE# or OE# change	^t OH	0		0		ns
READ cycle time	^t RC		70		80	ns

Figure 11: Output Load Circuit





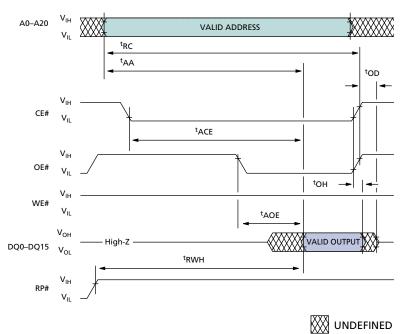
Write Cycle Timing Requirements

		-70)/-80	
PARAMETER	SYMBOL	MIN	MAX	UNITS
^t RST# HIGH recovery to WE# going LOW	^t RS	150		ns
CE# setup to WE# going LOW	tCS	0		ns
Write pulse width	^t WP	30		ns
Data setup to WE# going HIGH	^t DS	30		ns
Address setup to WE# HIGH	^t AS	35		ns
CE# hold from WE# HIGH	^t CH	0		ns
Data hold from WE# HIGH	^t DH	0		ns
Address hold from WE# HIGH	^t AH	0		ns
Write pulse width HIGH	tWPH	30		ns
WP# setup to WE# going HIGH	^t RHS	0		ns
VPP setup to WE# going HIGH	^t VPS	200		ns
Write recovery before READ	tWOS	50		ns
WP# hold from valid SRD	^t RHH	0		ns
VPP hold from valid SRD	^t VPH	0		ns
WE# HIGH to data valid	^t WB		^t AA + 50	ns

Erase and Program Cycle Timing Requirements

	-70/-80		
PARAMETER	ТҮР	MAX	UNITS
4KW parameter block program time	40	800	ms
32KW main block program time	320	6,400	ms
Word program time	8	10,000	μs
4KW parameter block erase time	0.3	6	S
32KW main block erase time	0.5	6	S
Program suspend latency	5	10	μs
Erase suspend latency	5	20	μs
Chip programming time (FPA)		20	S
FPA setup time	5		μs





Single Asynchronous Read Operation

Read Timing Parameters

	-7	0	-80		
SYMBOL	MIN	ΜΑΧ	MIN	МАХ	UNITS
^t AA		70		80	ns
^t ACE		70		80	ns
^t AOE		25		25	ns
^t RWH		200		200	ns
^t OD		25		25	ns
^t OH	0		0		ns
^t RC		70		80	ns



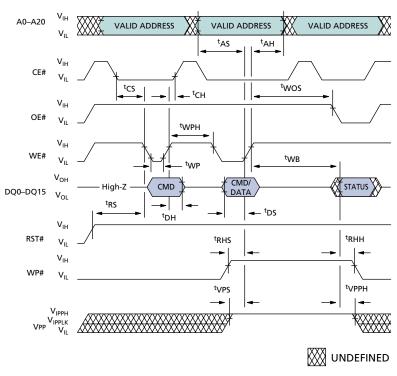
V_{IH} A3-A20 VALID ADDRESS VIL VIH VALID VALID VALID A0-A2 VALID ADDRESS VIL ^tAA tOD 1-VIH CE1# V_{IL} ^tACE VIH OE# VIL -{ } VIH WE# V_{IL} ^tAOE → 🖛 ^tAPA tОН VALIU V_{OH} VALID OUTPUT High-Z DQ0-DQ15 V_{OL} ^tRWH V_{IH} RP# VIL W UNDEFINED

Asynchronous Page Mode Read Operation

Read Timing Parameters

	-7	70	-80		
SYMBOL	MIN	ΜΑΧ	MIN	МАХ	UNITS
^t AA		70		80	ns
^t ACE		70		80	ns
^t APA		30		30	ns
^t AOE		25		25	ns
^t RWH		200		200	ns
^t OD		25		25	ns
^t OH	0		0		ns





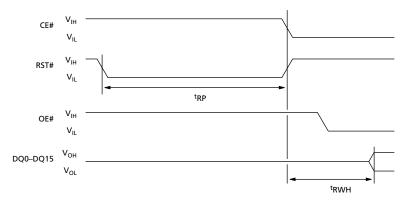
Two-Cycle Programming/Erase Operation

Read Timing Parameters

	-70		
SYMBOL	MIN	MAX	UNITS
^t RS	150		ns
^t CS	0		ns
^t WP	30		ns
^t DS	30		ns
^t AS	35		ns
^t CH	0		ns
^t DH	0		ns
^t AH	0		ns
^t RHS	0		ns
^t VPS	200		ns
tWOS	50		ns
^t RHH	0		ns
^t VPPH	0		ns
^t WB		^t AA + 50	ns



Reset Operation



Read and Write Timing Parameters

	-70/-80		
SYMBOL	MIN	ΜΑΧ	UNITS
^t RWH		200	ns
^t RP	100		ns



Table 12: CFI

OFFSET	DATA	DESCRIPTION
00	2Ch	Manufacturer code
01	94h	Top boot block device code
	95h	Bottom boot block device code
02ñ0F	reserved	Reserved
10, 11	0051, 0052	"QR"
12	0059	"Y"
13, 14	0003, 0000	Primary OEM command set
15, 16	0039, 0000	Address for primary extended table
17, 18	0000, 0000	Alternate OEM command set
19, 1A	0000, 0000	Address for OEM extended table
1B	0027	Vcc MIN for Erase/Write; Bit 7-bit 4 volts in BCD; Bit 3-bit 0 100mV in BCD
1C	0033	Vcc MAX for Erase/Write; Bit 7–bit 4 volts in BCD; Bit 3–bit 0 100mV in BCD
1D	00B4	VPP MIN for Erase/Write; Bit 7–bit 4 volts in hex; Bit 3–bit 0 100mV in BCD
1E	00C6	VPP MAX for Erase/Write; Bit 7–bit 4 volts in hex; Bit 3–bit 0 100mV in BCD
1F	0003	Typical timeout for single byte/word program, $2^{n}\mu s$, 0000 = not supported
20	0000	Typical timeout for maximum size multiple byte/word program, 2 ⁿ µs, 0000 = not supported
21	0009	Typical timeout for individual block erase, 2 ⁿ ms, 0000 = not supported
22	0000	Typical timeout for full chip erase, 2 ⁿ s, 0000 = not supported
23	000C	Maximum timeout for single byte/word program, $2^{n}\mu s$, 0000 = not supported
24	0000	Maximum timeout for maximum size multiple byte/word program, $2^n \mu s$, 0000 = not supported
25	0003	Maximum timeout for individual block erase, 2 ⁿ s, 0000 = not supported
26	0000	Maximum timeout for full chip erase, 2 ⁿ s, 0000 = not supported
27	0016	Device size, 2 ⁿ bytes
28	0001	Bus interface x16 = 1
29	0000	Flash device interface description 0000 = async
2A, 2B	0000, 0000	Maximum number of bytes in multi-byte program or page, 2 ⁿ
2C	0003	Number of erase block regions within device (4K words and 32K words)
2D, 2E	002F, 0000	Top boot block device erase block region information 1, 8 blocks
	0007, 0000	G
2F, 30	0000, 0001	Top boot block deviceof 8KB
	0020, 0000	Bottom boot block deviceof 8KB
31, 32	000E, 0000	•
	000E, 0000	Bottom boot block 15 blocks of
33, 34	0000, 0001	64КВ
35, 36	0007, 0000	Top boot block device48 blocks of
	002F, 0000	Bottom boot block device48 blocks of
37, 38	0020, 0000	Top boot block device64KB
	0000, 0001	Bottom boot block device64KB
39, 3A	0050, 0052	"PR"
3B	0049	
3C	0030	Major version number, ASCII

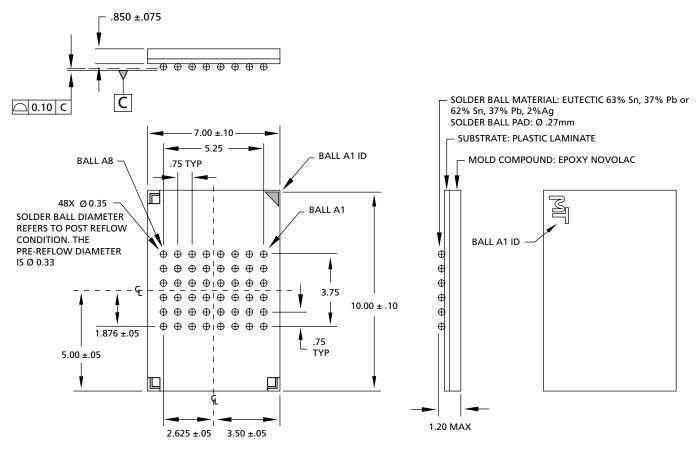


Table 12: CFI

OFFSET	DATA	DESCRIPTION
3D	0031	Minor version number, ASCII
3E	00E6	Optional Feature and Command Support
3F	0002	Bit 0 Chip erase supported no = 0
40	0000	Bit 1 Suspend erase supported = yes = 1
41	0000	Bit 2 Suspend program supported = yes = 1
		Bit 3 Chip lock/unlock supported = no = 0
		Bit 4 Queued erase supported = $no = 0$
		Bit 5 Instant individual block locking supported = yes = 1
		Bit 6 Protection bits supported = yes = 1
		Bit 7 Page mode read supported = yes =1
		Bit 8 Synchronous read supported = no =0 Bit 9 Simultaneous operation supported = yes = 1
42	0001	
42	0001	Program supported after erase suspend = yes
43, 44	0003, 0000	Bit 0 block lock status active = yes; Bit 1 block lock down active = yes
45	0030	Vcc supply optimum, 00 = not supported, Bit 7-bit 4 volts in BCD; Bit 3-bit 0 100mV in BCD
46	00C0	VPP supply optimum, 00 = not supported, Bit 7-bit 4 volts in BCD; Bit 3-bit 0 100mV in BCD
47	0001	Number of protection register fields in JEDEC ID space
48, 49		Lock bytes LOW address, lock bytes HIGH address
4A, 4B	0003, 0003	2 ⁿ factory programmed bytes, 2 ⁿ user programmable bytes
4C	0003	Background Operation
		0000= Not used
		0001 = 4% block split
		0002 = 12% block split
		0003 = 25% block split
4D	0000	0004 = 505 block split
40	0000	Burst Mode Type 0000 = No burst mode
		00x1 = 4 words MAX
		00x2 = 8 words MAX
		00x3 = 16 words MAX
		001x = Linear burst, and/or
		002x = interleaved burst, and/or
		004x = Continuous burst
4E	0002	Page Mode Type
		0000 = No page mode
		0001 = 4-word page
		0002 = 8-word page
		0003 = 16-word page
		0004 = 32-word page
4F	0000	Not used



48-Ball FBGA



NOTE:

1. All dimensions in millimeters.

2. Package width and length do not include mold protrusion; allowable mold protrusion is 0.27mm per side.

DATA SHEET DESIGNATION

No Mark: This data sheet contains minimum and maximum limits specified over the complete power supply and temperature range for production devices. Although considered final, these specifications are subject to change, as further product development and data characterization sometimes occur.



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ev.F
 Clarified column heading on Table 6 Clarified default condition after PROGRAM command Clarified default condition after ERASE command Clarified DQ0 value for PR_LOCK register Updated value for ^tAS
ev. E
 Added Programming Chip Protection Register section Updated Automated Word Programming and Block Erase flowcharts Updated Read Query, Programming Operations, Erase Operations, Unlocked State, and Read-While-Write/
Erase Concurrency sectionsAdded updated version of the Command State Machine Transition State table
ev. 4
 ev. 3, PRELIMINARY7/02 Updated Status Register section Updated command descriptions Updated flowcharts Updated Read-While-Write/Erase Concurrency section Updated timing diagrams
ev. 2, PRELIMINARY
 Changed ^tAH from 90ns to 0ns Updated Read Chip Protection Identification Data text
 ev. 2, PRELIMINARY
Updated timing diagrams and parameters priginal document, PRELIMINARY